MULTINOMIAL LOGISTIC REGRESSION ALGORITHM* **

DANKMAR BÖHNING

Department of Epidemiology, Free University Berlin, Augustastr. 37
1000 Berlin 45, Germany

(Received July 23, 1990; revised October 12, 1990)

Abstract. The lower bound principle (introduced in Böhnning and Lindsay (1988, Ann. Inst. Statist. Math., 40, 641–663), Böhnning (1989, Biometrika, 76, 375–383) consists of replacing the second derivative matrix by a global lower bound in the Loewner ordering. This bound is used in the Newton-Raphson iteration instead of the Hessian matrix leading to a monotonically converging sequence of iterates. Here, we apply this principle to the multinomial logistic regression model, where it becomes specifically attractive.

Key words and phrases: Kronecker product, Loewner ordering, lower bound principle, monotonicity.

1. Introduction

Let $L(\pi)$ denote the log-likelihood, $\nabla L(\pi)$ the score vector and $\nabla^2 L(\pi)$ the second derivative matrix at $\pi \in \mathbb{R}^m$. Suppose

$$\nabla^2 L(\pi) \geq B$$

for all $\pi$ and some negative definite $m \times m$ matrix $B$. Here $C \geq D$ denotes Loewner ordering of two matrices and means that $C - D$ is non-negative definite. Consider the second order Taylor series for the log-likelihood at $\pi_0$:

$$L(\pi) - L(\pi_0) = (\pi - \pi_0)^T \nabla L(\pi_0) + \frac{1}{2}(\pi - \pi_0)^T \nabla^2 L(\pi_0 + \alpha(\pi - \pi_0))(\pi - \pi_0)$$

$$\geq (\pi - \pi_0)^T \nabla L(\pi_0) + \frac{1}{2}(\pi - \pi_0)^T B(\pi - \pi_0)$$

where we have used (1.1) to achieve the lower bound for $L$. Maximizing the right-hand side of the above inequality yields the Lower Bound iterate $\pi_{LB} = \pi_0 - B^{-1} \nabla L(\pi_0)$. We have the following:


** This research was supported by the German Research Foundation.
Theorem 1.1. (Böhning and Lindsay (1988))

(i) (Monotonicity) For the Lower Bound iterate we have

\[ L(\pi_{\text{LB}}) \geq L(\pi_0) \quad \text{with} \quad " > " \quad \text{if} \quad \pi_{\text{LB}} \neq \pi_0. \]

(ii) (Convergence) Let \((\pi_j)\) be a sequence created by the lower bound algorithm. If \(L\) is bounded above in addition, then

\[ \|\nabla L(\pi_j)\| \to 0. \]

2. Multinomial logistic regression

We observe vectors \(Y = (y_1, \ldots, y_{k+1})^T\), with \(y_i = 0\) for all \(i\) besides one \(j\) with \(y_j = 1\) and corresponding probability \(p_j\), implying

\[ EY = p, \quad \text{Cov} \, Y = \Lambda_p - pp^T, \quad \Lambda_p = \begin{pmatrix} p_1 & \cdots & 0 \\ 0 & \ddots & \vdots \\ \vdots & \ddots & p_{k+1} \end{pmatrix}. \]

Recall that the multinomial logit-model is given by

\[ p_i = \frac{\exp(\pi^{(i)^T} x)}{1 + \sum_{j=1}^{k} \exp(\pi^{(j)^T} x)} \quad \text{for} \quad i = 1, \ldots, k, \]

\[ p_{k+1} = 1 \frac{1}{1 + \sum_{j=1}^{k} \exp(\pi^{(j)^T} x)} \]

where \(x = (x_1, \ldots, x_m)^T\) is the vector of covariates, and \(\pi^{(i)}\) is the parameter vector corresponding to the \(i\)-th response category. For reasons of simplicity in presentation, consider the log-likelihood of just one observation \(Y\):

\[ \log \prod_{j=1}^{k+1} p_j^{y_j} = \sum_{j=1}^{k} y_j \pi^{(j)^T} x - \log \left[ 1 + \sum_{j=1}^{k} \exp(\pi^{(j)^T} x) \right]. \]

Let \(\pi = (\pi_1^{(1)}, \ldots, \pi_m^{(1)}, \ldots, \pi_1^{(k)}, \ldots, \pi_m^{(k)})^T\) denote the \(mk\)-vector of \(mk\) parameters, the upper index going along with the response category, the lower index with the covariate. We have for the partial derivative

\[ \frac{\partial L}{\partial \pi^{(h)}_g} = y_h x_g - \frac{\exp(\pi^{(h)^T} x)}{1 + \sum_{j=1}^{k} \exp(\pi^{(j)^T} x)} x_g = (y_h - \hat{p}_h)x_g \]

with the notation \(\hat{p}_h = \frac{\exp(\pi^{(h)^T} x)}{1 + \sum_{j=1}^{k} \exp(\pi^{(j)^T} x)}\). This yields the score vector

\[ \nabla L(\pi) = [(y_1 - \hat{p}_1)x_1, \ldots, (y_1 - \hat{p}_1)x_m, \ldots, (y_k - \hat{p}_k)x_1, \ldots, (y_k - \hat{p}_k)x_m]^T = (Y - \hat{p}) \otimes x \]
where $\otimes$ is the Kronecker product $A \otimes B$ of two arbitrary matrices. The observed information can be easily computed to be

$$
\frac{\partial^2 L}{\partial \pi^{(h')} \partial \pi^{(h)}} = \delta_{hh'} \exp(\pi^{(h')}^T x) \left( 1 + \sum_{j=1}^k \exp(\pi^{(j)}^T x) \right) - \exp(\pi^{(h')}^T x) \exp(\pi^{(h)}^T x) \right) \right) x_g^T x_g
$$

leading to the observed information matrix

$$
-\nabla^2 L = \begin{pmatrix}
\hat{p}_1(1 - \hat{p}_1) xx^T & -\hat{p}_1 \hat{p}_2 xx^T & \cdots & -\hat{p}_1 \hat{p}_k xx^T \\
-\hat{p}_2 \hat{p}_1 xx^T & \hat{p}_2(1 - \hat{p}_2) xx^T & \cdots & \hat{p}_2 \hat{p}_k xx^T \\
& \cdots & \cdots & \cdots \\
-\hat{p}_k \hat{p}_1 xx^T & \cdots & \hat{p}_k(1 - \hat{p}_k) xx^T & \end{pmatrix}
$$

The proof of the following lemma is straightforward.

**Lemma 2.1.** If $A \preceq B$ then for symmetric, nonnegative definite $C$:

$$
A \otimes C \preceq B \otimes C.
$$

**Lemma 2.2.** $\Lambda_p - pp^T \leq [E - 11^T/(k+1)]/2$, where $1$ is the $k$-vector of $1$'s.

A proof of this lemma is given in the proof of Theorem 5.3 in Böhning and Lindsay (1988) or can be constructed from Theorem 2 in Baksalary and Pukelsheim (1985).

Taking Lemmas 2.1 and 2.2 together, we get the following main result:

**Theorem 2.1.** (a) For the information matrix of one observation we have:

$$
i(\pi) = (\Lambda_\pi - \hat{p}\hat{p}^T) \otimes xx^T \leq \frac{1}{2} [E - 11^T/(k+1)] \otimes xx^T.
$$

(b) For the information matrix of a sample of size $n$ we get:

$$
i_{\text{com}}(\pi) = \sum_{i=1}^n (\Lambda_{\hat{p}_i} - \hat{p}_i \hat{p}_i^T) \otimes x_i x_i^T \leq \sum_{i=1}^n \frac{1}{2} [E - 11^T/(k+1)] \otimes x_i x_i^T
$$

$$
= \frac{1}{2} [E - 11^T/(k+1)] \otimes \sum_{i=1}^n x_i x_i^T
$$

$$
= \frac{1}{2} [E - 11^T/(k+1)] \otimes X^T X =: B.
$$
where $X$ is the $n \times m$ design matrix
\[
\begin{pmatrix}
  x_1^T \\
  \vdots \\
  x_n^T
\end{pmatrix}.
\]

(c) $B^{-1} = 2[E - 11^T/(k + 1)]^{-1} \otimes (X^T X)^{-1} = 2[E + 11^T] \otimes (X^T X)^{-1}.$

(d) $\pi_{LB} = \pi_0 + B^{-1} \sum_{i=1}^n (y_i - \hat{p}_i) \otimes x_i.$

**Remark.** Since $\sum_{i=1}^n (\Lambda \tilde{p}_i - \hat{p}_i \hat{p}_i^T) \otimes x_i x_i^T = \sum_{i=1}^n (\Lambda \tilde{p}_i - \hat{p}_i \hat{p}_i^T) \otimes \sum_{i=1}^n x_i x_i^T$ is not true in general, we would have to invert the $(mk)^2$ matrix $i_{com}$ at each step of the Newton-Raphson iteration. If we have 6 response categories ($k = 5$) and $m = 10$ covariates, then $i_{com}$ is a $50 \times 50$ matrix. In contrast, the lower bound principle needs to invert a $10 \times 10$ matrix only once. The lower bound algorithm converges linearly with convergence rate depending on $\|E - B^{-1} V_2 L(\hat{\pi})\|$. If $\hat{\pi} = 0$, then the lower bound algorithm converges at least superlinearly. Thus, if $\hat{\pi}$ is "near" zero, the computational efficiency of the lower bound iteration can be expected to be better than that of the Newton-Raphson iteration. To evaluate this point, in Böning and Lindsay (1988), Section 5.1) a simulation experiment was undertaken for binomial logistic regression, that is $k = 1$. There, the comparison is essentially between inverting a $k \times k$ matrix once (the lower bound algorithm) and inverting it several times (until a stopping rule is met, for the Newton-Raphson iteration). In all cases studied there, the computational efficiency of the lower bound method was better than that of the Newton-Raphson iteration. However, a downward-tendency was observed when the difference in CPU-time was plotted against distance of $\hat{\pi}$ to zero. Thus, it is possible that in extreme cases the Newton-Raphson algorithm might be more efficient. Here, we are comparing the single inversion of a $k \times k$ matrix (in the lower bound algorithm) with several inversions of a $km \times km$ matrix (in the Newton-Raphson iteration). This feature makes the lower bound method specifically attractive.

Acknowledgement

The author thanks an unknown referee for helpful comments and improvements.

**References**

