MULTINOMIAL LOGISTIC REGRESSION ALGORITHM* **

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Abstract. The lower bound principle (introduced in Böhning and Lindsay (1988, Ann. Inst. Statist. Math., 40, 641–663), Böhning (1989, Biometrika, 76, 375–383) consists of replacing the second derivative matrix by a global lower bound in the Loewner ordering. This bound is used in the Newton-Raphson iteration instead of the Hessian matrix leading to a monotonically converging sequence of iterates. Here, we apply this principle to the multinomial logistic regression model, where it becomes specifically attractive.

Key words and phrases: Kronecker product, Loewner ordering, lower bound principle, monotonicity.

1. Introduction

Let $L(\pi)$ denote the log-likelihood, $\nabla L(\pi)$ the score vector and $\nabla^2 L(\pi)$ the second derivative matrix at $\pi \in \mathbb{R}_m$. Suppose

(1.1)
$$\nabla^2 L(\pi) \ge B$$

for all π and some negative definite $m \times m$ matrix B. Here $C \ge D$ denotes Loewner ordering of two matrices and means that C-D is non-negative definite. Consider the second order Taylor series for the log-likelihood at π_0 :

$$L(\pi) - L(\pi_0) = (\pi - \pi_0)^T \nabla L(\pi_0) + \frac{1}{2} (\pi - \pi_0)^T \nabla^2 L(\pi_0 + \alpha (\pi - \pi_0))(\pi - \pi_0)$$

$$\geq (\pi - \pi_0)^T \nabla L(\pi_0) + \frac{1}{2} (\pi - \pi_0)^T B(\pi - \pi_0)$$

where we have used (1.1) to achieve the lower bound for L. Maximizing the right-hand side of the above inequality yields the Lower Bound iterate $\pi_{\rm LB} = \pi_0 - B^{-1} \nabla L(\pi_0)$. We have the following:

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THEOREM 1.1. (Böhning and Lindsay (1988)) (i) (Monotonicity) For the Lower Bound iterate we have

$$L(\pi_{\text{LB}}) \ge L(\pi_0)$$
 with ">" if $\pi_{\text{LB}} \ne \pi_0$

(ii) (Convergence) Let (π_j) be a sequence created by the lower bound algorithm. If L is bounded above in addition, then

$$\left\|\nabla L(\pi_j)\right\| \xrightarrow[j \to \infty]{} 0$$

2. Multinomial logistic regression

We observe vectors $Y = (y_1, \ldots, y_{k+1})^T$, with $y_i = 0$ for all *i* besides one *j* with $y_j = 1$ and corresponding probability p_j , implying

$$\boldsymbol{E}Y = \boldsymbol{p}, \quad \operatorname{Cov} Y = \Lambda_p - pp^T, \quad \Lambda_p = \begin{pmatrix} p_1 & 0\\ 0 & \cdots & 0\\ \vdots & \vdots\\ 0 & p_{k+1} \end{pmatrix}.$$

Recall that the *multinomial logit-model* is given by

$$p_{i} = \exp(\pi^{(i)^{T}} \boldsymbol{x}) / \left[1 + \sum_{j=1}^{k} \exp(\pi^{(j)^{T}} \boldsymbol{x}) \right] \quad \text{for} \quad i = 1, \dots, k,$$
$$p_{k+1} = 1 / \left[1 + \sum_{j=1}^{k} \exp(\pi^{(j)^{T}} \boldsymbol{x}) \right]$$

where $\boldsymbol{x} = (x_1, \ldots, x_m)^T$ is the vector of covariates, and $\pi^{(i)}$ is the parameter vector corresponding to the *i*-th response category. For reasons of simplicity in presentation, consider the log-likelihood of just *one* observation Y:

$$\log \prod_{j=1}^{k+1} p_j^{y_j} = \sum_{j=1}^k y_j \pi^{(j)^T} \boldsymbol{x} - \log \left[1 + \sum_{j=1}^k \exp(\pi^{(j)^T} \boldsymbol{x}) \right].$$

Let $\pi = (\pi_1^{(1)}, \ldots, \pi_m^{(1)}, \ldots, \pi_1^{(k)}, \ldots, \pi_m^{(k)})^T$ denote the *mk*-vector of *mk* parameters, the upper index going along with the response category, the lower index with the covariate. We have for the partial derivative

$$\frac{\partial L}{\partial \pi_g^{(h)}} = y_h x_g - \frac{\exp(\pi^{(h)^T} x)}{1 + \sum_{j=1}^k \exp(\pi^{(j)^T} x)} x_g = (y_h - \hat{p}_h) x_g$$

with the notation $\hat{p}_h = \exp(\pi^{(h)^T} \boldsymbol{x}) / (1 + \sum_{j=1}^k \exp(\pi^{(j)^T} \boldsymbol{x}))$. This yields the score vector

$$egin{aligned}
abla L(\pi) &= [(y_1 - \hat{p}_1) x_1, \dots, (y_1 - \hat{p}_1) x_m, \dots, (y_k - \hat{p}_k) x_1, \dots, (y_k - \hat{p}_k) x_m]^T \ &= (Y - \hat{oldsymbol{p}}) \otimes oldsymbol{x} \end{aligned}$$

where \otimes is the Kronecker product $A \otimes B$ of two arbitrary matrices. The observed information can be easily computed to be

$$-\frac{\partial^{2} L}{\partial \pi_{g'}^{(h')} \partial \pi_{g}^{(h)}} = \frac{\delta_{hh'} \exp(\pi^{(h)^{T}} \boldsymbol{x}) \left(1 + \sum_{j=1}^{k} \exp(\pi^{(j)^{T}} \boldsymbol{x})\right) - \exp(\pi^{(h')^{T}} \boldsymbol{x}) \exp(\pi^{(h)^{T}} \boldsymbol{x})}{\left(1 + \sum_{j=1}^{k} \exp(\pi^{(j)^{T}} \boldsymbol{x})\right)^{2}} x_{g'} x_{g}$$

$$= (\delta_{hh'} \hat{p}_{h} - \hat{p}_{h'} \hat{p}_{h}) x_{g'} x_{g},$$

leading to the observed information matrix

$$-\nabla^2 L = \begin{pmatrix} \hat{p}_1(1-\hat{p}_1)\boldsymbol{x}\boldsymbol{x}^T & -\hat{p}_1\hat{p}_2\boldsymbol{x}\boldsymbol{x}^T & \cdots & -\hat{p}_1\hat{p}_k\boldsymbol{x}\boldsymbol{x}^T \\ \vdots & \hat{p}_2(1-\hat{p}_2)\boldsymbol{x}\boldsymbol{x}^T & \vdots \\ -\hat{p}_k\hat{p}_1\boldsymbol{x}\boldsymbol{x}^T & \cdots & \cdots & \hat{p}_k(1-\hat{p}_k)\boldsymbol{x}\boldsymbol{x}^T \end{pmatrix}$$
$$= (\Lambda_{\hat{p}} - \hat{p}\hat{p}^T) \otimes \boldsymbol{x}\boldsymbol{x}^T.$$

The proof of the following lemma is straightforward.

LEMMA 2.1. If $A \leq B$ then for symmetric, nonnegative definite C:

$$A \otimes C \leq B \otimes C.$$

LEMMA 2.2. $\Lambda_p - pp^T \leq [E - \mathbf{11}^T/(k+1)]/2$, where **1** is the k-vector of 1's.

A proof of this lemma is given in the proof of Theorem 5.3 in Böhning and Lindsay (1988) or can be constructed from Theorem 2 in Baksalary and Pukelsheim (1985).

Taking Lemmas 2.1 and 2.2 together, we get the following main result:

THEOREM 2.1. (a) For the information matrix of one observation we have:

$$i(\pi) = (\Lambda_{\hat{p}} - \hat{p}\hat{p}^T) \otimes \boldsymbol{x}\boldsymbol{x}^T \leq \frac{1}{2}[E - \mathbf{1}\mathbf{1}^T/(k+1)] \otimes \boldsymbol{x}\boldsymbol{x}^T.$$

(b) For the information matrix of a **sample** of size n we get:

$$\begin{split} i_{\text{com}}(\pi) &= \sum_{i=1}^{n} (\Lambda_{\hat{p}_{i}} - \hat{p}_{i} \hat{p}_{i}^{T}) \otimes \mathbf{x}_{i} \mathbf{x}_{i}^{T} \leq \sum_{i=1}^{n} \frac{1}{2} [E - \mathbf{1} \mathbf{1}^{T} / (k+1)] \otimes \mathbf{x}_{i} \mathbf{x}_{i}^{T} \\ &= \frac{1}{2} [E - \mathbf{1} \mathbf{1}^{T} / (k+1)] \otimes \sum_{i=1}^{n} \mathbf{x}_{i} \mathbf{x}_{i}^{T} \\ &= \frac{1}{2} [E - \mathbf{1} \mathbf{1}^{T} / (k+1)] \otimes X^{T} X =: B, \end{split}$$

where X is the $n \times m$ design matrix $\begin{pmatrix} \boldsymbol{x}_1^T \\ \vdots \\ \ddots \end{pmatrix}$.

$$\begin{array}{c} \left\langle \boldsymbol{x}_{n}^{T} \right\rangle \\ \text{(c)} \ B^{-1} = 2[E - \mathbf{1}\mathbf{1}^{T}/(k+1)]^{-1} \otimes (X^{T}X)^{-1} = 2[E + \mathbf{1}\mathbf{1}^{T}] \otimes (X^{T}X)^{-1} \\ \text{(d)} \ \pi_{\text{LB}} = \pi_{0} + B^{-1} \sum_{i=1}^{n} (\boldsymbol{Y}_{i} - \hat{\boldsymbol{p}}_{i}) \otimes \boldsymbol{x}_{i}. \end{array}$$

Remark. Since $\sum_{i=1}^{n} (\Lambda_{\hat{p}_{i}} - \hat{p}_{i} \hat{p}_{i}^{T}) \otimes \boldsymbol{x}_{i} \boldsymbol{x}_{i}^{T} = \sum_{i=1}^{n} (\Lambda_{\hat{p}_{i}} - \hat{p}_{i} \hat{p}_{i}^{T}) \otimes \sum_{i=1}^{n} \boldsymbol{x}_{i} \boldsymbol{x}_{i}^{T}$ is not true in general, we would have to invert the $(mk)^{2}$ matrix i_{com} at each step of the Newton-Raphson iteration. If we have 6 response categories (k = 5) and m = 10 covariates, then $i_{\rm com}$ is a 50 \times 50 matrix. In contrast, the lower bound principle needs to invert a 10×10 matrix only once. The lower bound algorithm converges linearly with convergence rate depending on $||E - B^{-1}\nabla^2 L(\hat{\pi})||$. If $\hat{\pi} = 0$, then the lower bound algorithm converges at least superlinearly. Thus, if $\hat{\pi}$ is "near" zero, the computational efficiency of the lower bound iteration can be expected to be better than that of the Newton-Raphson iteration. To evaluate this point, in Böhning and Lindsay ((1988), Section 5.1) a simulation experiment was undertaken for *binomial* logistic regression, that is k = 1. There, the comparison is essentially between inverting a $k \times k$ matrix once (the lower bound algorithm) and inverting it several times (until a stopping rule is met, for the Newton-Raphson iteration). In all cases studied there, the computational efficiency of the lower bound method was better than that of the Newton-Raphson iteration. However, a downward-tendency was observed when the difference in CPU-time was plotted against distance of $\hat{\pi}$ to zero. Thus, it is possible that in extreme cases the Newton-Raphson algorithm might be more efficient. Here, we are comparing the single inversion of a $k \times k$ matrix (in the lower bound algorithm) with several inversions of a $km \times km$ matrix (in the Newton-Raphson iteration). This feature makes the lower bound method specifically attractive.

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